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Alphabet-almost-simple 2-neighbour-transitive codes

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Abstract

Let X be a subgroup of the full automorphism group of the Hamming graph $H(m, q)$, and C a subset of the vertices of the Hamming graph. We say that C is an $(X, 2)$ -neighbour-transitive code if X is transitive on C , as well as C_1 and C_2 , the sets of vertices which are distance 1 and 2 from the code. It has been shown that, given an $(X, 2)$ -neighbour-transitive code C , there exists a subgroup of X with a 2-transitive action on the alphabet; this action is thus almost-simple or affine. This paper completes the classification of $(X, 2)$ -neighbour-transitive codes, with minimum distance at least 5, where the subgroup of X stabilising some entry has an almost-simple action on the alphabet in the stabilised entry. The main result of this paper states that the class of $(X, 2)$ neighbour-transitive codes with an almost-simple action on the alphabet and minimum distance at least 3 consists of one infinite family of well known codes.

Keywords: 2-neighbour-transitive, alphabet-almost-simple, automorphism groups, Hamming graph, completely transitive.

Math. Subj. Class.: 05E20, 68R05, 20B25

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1 Introduction

Ever since Shannon's 1948 paper [18, 19] there has been a great deal of interest around families of error-correcting codes with a high degree of symmetry. The rationale behind this interest is that codes with symmetry should have good error correcting properties. The first families classified were perfect (see [21] or [22]) and nearly-perfect (defined in [12], classified in [15]) codes over prime power alphabets. Note that the classification of nearly-perfect codes follows from the earlier results of [17] on uniformly packed codes, since nearly-perfect codes are uniformly packed codes with maximal packing density. These classifications show that perfect and nearly-perfect codes are rare. In an effort to find further classes of efficient codes, Delsarte [4] introduced *completely regular* codes, a more general class of codes that possess a high degree of combinatorial symmetry. Much effort has been put into classifying particular classes of completely regular codes (see for instance [1, 2]), and new completely regular codes continue to be found [6]. However, completely regular codes have proven to be hard to classify, and this remains an open problem.

Completely transitive codes (first defined in [20], with a generalisation studied in [10]) are a class of codes with a high degree of algebraic symmetry and are a subset of completely regular codes. As such a classification of completely transitive codes would be interesting from the point of view of classifying completely regular codes. This problem also remains open.

Here, the conditions of complete transitivity are relaxed and the family of *2-neighbour-transitive codes* is studied, a class of codes with a moderate degree of algebraic symmetry. Note that every completely transitive code (see Section 2) is 2-neighbour-transitive. By studying this class of codes we hope to find new codes and gain a better understanding of completely transitive codes. Indeed a classification of 2-neighbour-transitive codes would have as a corollary a classification of completely transitive codes. We also note that codes with 2-transitive actions on the entries of the Hamming graph (which 2-neighbour-transitive codes indeed have), have been of interest lately, where this fact can be used to prove that certain families of codes achieve capacity on erasure channels [14]. The analysis of 2-neighbour-transitive codes is being attacked as three separate problems: *entry-faithful* (see [7]), *alphabet-almost-simple*, and *alphabet-affine*. This paper concerns the alphabet-almost-simple case. The results of this paper do not return any new examples.

However, the results here are of interest from the point of view of perfect codes over an alphabet of non-prime-power size, since in this case a code cannot be alphabet-affine (and also not entry-faithful, by [7]), but may be alphabet-almost-simple. The existence of perfect codes over non-prime-power alphabets with covering radius 1 or 2, is still an open question (see [13]). By Theorem 1.1, if such codes exist, then they cannot be 2-neighbour-transitive (unless they are equivalent to the repetition code of length 3). Note that in the prime power case, for each set of parameters for which a perfect code with covering radius $\rho \geq 2$ exists, a 2-neighbour-transitive code with those parameters exists. That is, the repetition and Golay codes are 2-neighbour-transitive. In fact, the repetition, Hamming and Golay codes are completely transitive (by [11, Example 3.1] for the repetition codes, [20, Proposition 7.3] for the Hamming and binary Golay codes, and [10, Example 3.5.6] for the ternary Golay codes).

1.1 Statement of the main results

Let X be a subgroup of the full automorphism group $S_q^m \rtimes S_m$ of the Hamming graph $\Gamma = H(m, q)$ and let C be a code, that is, a subset of the set of vertices $V\Gamma$. We say that C is an (X, s) -neighbour-transitive code if X fixes C setwise and acts transitively on $C = C_0, C_1, \dots, C_s$ (where C_i are parts of the distance partition, see Section 2). In joint work with Giudici and Praeger [7], the authors classified all $(X, 2)$ -neighbour-transitive codes for which the group X acts faithfully on the set of entries of the Hamming graph. In this paper, we begin the study of $(X, 2)$ -neighbour-transitive codes such that the action of X on the entries has a non-trivial kernel.

Let M be the set of entries of the Hamming graph $H(m, q)$ and Q_i be the copy of the alphabet Q in the corresponding entry $i \in M$. Then the vertex set of $H(m, q)$ is:

$$V\Gamma = \prod_{i \in M} Q_i.$$

If C is an $(X, 2)$ -neighbour-transitive code with minimum distance $\delta \geq 3$, then the subgroup $X_i \leq X$ stabilising the entry $i \in M$ has a 2-transitive action on the alphabet Q_i in that entry (see [7, Proposition 2.7]). Any 2-transitive group G is of affine type ($G \leq \text{AGL}_d(p)$ for some integer d and prime p) or almost-simple type ($S \leq G \leq \text{Aut}(S)$ for some non-abelian simple group S) [5, Theorem 4.1B]. If the action of X on M (see Section 2.1) is transitive with a non-trivial kernel and the action of X_i on the alphabet Q_i is almost-simple then we say C is X -alphabet-almost-simple. Our main aim here is to prove the non-existence of codes which are X -alphabet-almost-simple and $(X, 2)$ -neighbour-transitive with minimum distance $\delta \geq 4$.

Theorem 1.1. *Let C be an X -alphabet-almost-simple and $(X, 2)$ -neighbour-transitive code in $H(m, q)$ with minimum distance $\delta \geq 3$. Then $\delta = 3$ and C is equivalent to the repetition code in $H(3, q)$, where $q \geq 5$.*

In Section 2 we define the notation used in the paper. In Section 3 we give some results on the structure of codes that are X -alphabet-almost-simple and $(X, 2)$ -neighbour-transitive, as well as pose some questions about codes for which the action of X_i on the alphabet in the entry $i \in M$ is affine. We present some examples of codes with properties of interest in relation to our results in Section 4. Finally, in Section 5, we give a classification of diagonally $(X, 2)$ -neighbour-transitive codes (see Definition 3.1) and prove Theorem 1.1.

2 Preliminaries

Throughout this paper we let $M = \{1, \dots, m\}$ and $Q = \{1, \dots, q\}$, with $m, q \geq 2$, though if $q = 2$ we will at times use $Q = \{0, 1\}$. We refer to M as the *set of entries* and Q as the *alphabet*. We use Q_i to denote the disjoint copy of the alphabet Q in the entry $i \in M$. The vertex set $V\Gamma$ of the Hamming graph $\Gamma = H(m, q)$ consists of all m -tuples with entries labeled by the set M , taken from the set Q . An edge exists between two vertices if they differ as m -tuples in exactly one entry. For vertices α, β of $H(m, q)$ the *Hamming distance* $d(\alpha, \beta)$ is the number of entries in which α and β differ, i.e. the usual graph distance in Γ . For $\alpha \in V\Gamma$, we refer to the element of Q appearing in the i -th entry of α as α_i , so that $\alpha = (\alpha_1, \dots, \alpha_m)$ throughout.

A code C is a subset of the vertex set of the Hamming graph. The *minimum distance* of C is $\delta = \min\{d(\alpha, \beta) \mid \alpha, \beta \in C, \alpha \neq \beta\}$. For a vertex $\alpha \in H(m, q)$, define

$$\Gamma_r(\alpha) = \{\beta \in \Gamma \mid d(\alpha, \beta) = r\}, \quad \text{and} \quad d(\alpha, C) = \min\{d(\alpha, \beta) \mid \beta \in C\}.$$

We then define the *covering radius* to be

$$\rho = \max\{d(\alpha, C) \mid \alpha \in \Gamma\}.$$

For any $r \leq \rho$, define $C_r = \{\alpha \in \Gamma \mid d(\alpha, C) = r\}$. Note that C_i is the disjoint union $\cup_{\alpha \in C} \Gamma_i(\alpha)$ for $i \leq \lfloor \frac{\delta-1}{2} \rfloor$.

2.1 Automorphism groups

The automorphism group $\text{Aut}(\Gamma)$ of the Hamming graph is the semi-direct product $B \rtimes L$, where $B \cong S_q^m$ and $L \cong S_m$ (see [3, Theorem 9.2.1]). We refer to B as the *base group*, and L as the *top group*, of $\text{Aut}(\Gamma)$. Let $g = (g_1, \dots, g_m) \in B$, $\sigma \in L$ and α be a vertex in $H(m, q)$. Then g and σ act on α as follows:

$$\alpha^g = (\alpha_1^{g_1}, \dots, \alpha_m^{g_m}) \quad \text{and} \quad \alpha^\sigma = (\alpha_{1\sigma^{-1}}, \dots, \alpha_{m\sigma^{-1}}).$$

We define the automorphism group of a code C in $H(m, q)$ to be $\text{Aut}(C) = \text{Aut}(\Gamma)_C$, the setwise stabiliser of C in $\text{Aut}(\Gamma)$. For a subgroup $X \leq \text{Aut}(C)$ we define two other important actions of X which will be useful to us. First, consider the action of X on the set of entries M , which we will write as X^M . In particular $X^M = \mu(X)$, that is, the image of the homomorphism:

$$\mu : \begin{array}{ccc} X & \longrightarrow & S_m \\ (h_1, \dots, h_m)\sigma & \longmapsto & \sigma \end{array} .$$

Note that σ here is not necessarily an automorphism of C , that is, σ is a permutation of M but may not necessarily fix C setwise, though its pre-image $(h_1, \dots, h_m)\sigma$ is an element of $\text{Aut}(C)$. We define K to be the kernel of the map μ and note that $K = X \cap B$. In this paper we are concerned with $(X, 2)$ -neighbour-transitive codes where $K \neq 1$.

We also consider the action of the stabiliser $X_i \leq X$ of the entry $i \in M$, on the alphabet Q_i in that entry. We denote this action by $X_i^{Q_i} = \varphi_i(X_i)$, and it is the image of the homomorphism:

$$\varphi_i : \begin{array}{ccc} X_i & \longrightarrow & S_q \\ (h_1, \dots, h_m)\sigma & \longmapsto & h_i \end{array} .$$

Let C be a code in $H(m, q)$ and let X be a subgroup of $\text{Aut}(\Gamma)$. Recall that C is (X, s) -neighbour-transitive if each C_i is an X -orbit for $i = 0, \dots, s$. Note that this implies $X \leq \text{Aut}(C)$ and C is also (X, r) -neighbour-transitive, for $r < s$. If $s = 1$ then C is simply X -neighbour-transitive and if $s = \rho$, the covering radius, then C is X -completely transitive.

An *almost-simple* group is a group G where $S \leq G \leq \text{Aut}(S)$ for some non-abelian simple group S . The *socle* of a group G , denoted $\text{soc}(G)$, is the product of its minimal normal subgroups. The socle of an almost-simple group G is the non-abelian simple group S such that $S \leq G \leq \text{Aut}(S)$. Recall, if C is a code and $X \leq \text{Aut}(C)$ such that $K \neq 1$,

X^M is transitive on M and the $X_i^{Q_i}$ is almost-simple, then we say C is X -alphabet-almost-simple. We may sometimes omit the group X from any of the above terms, if the meaning is clear from the context.

We say that two codes, C and C' , in $H(m, q)$, are *equivalent* if there exists $x \in \text{Aut}(\Gamma)$ such that $C^x = C'$. Since elements of $\text{Aut}(\Gamma)$ preserve distance, equivalence preserves minimum distance.

2.2 Projections

For a subset $J = \{j_1, \dots, j_k\} \subseteq M$ we define the *projection of α with respect to J* as $\pi_J(\alpha) = (\alpha_{j_1}, \dots, \alpha_{j_k})$. For a code C we then define the *projection of C with respect to J* as $\pi_J(C) = \{\pi_J(\alpha) \mid \alpha \in C\}$. So π_J maps a vertex or code from $H(m, q)$ into the smaller Hamming graph $H(k, q)$.

Let X_J be the setwise stabiliser of a subset $J = \{j_1, \dots, j_k\} \subseteq M$. For $x = (h_1, \dots, h_m)\sigma \in X_J$, we define the *projection of x with respect to J* as $\chi_J(x)$ where

$$\pi_J(\alpha)^{\chi_J(x)} = \pi_J(\alpha^x).$$

To be well defined, this requires $x \in X_J$ and it follows that

$$\chi_J(x) = (h_{j_1}, \dots, h_{j_k})\hat{\sigma} \in \text{Aut}(H(k, q)),$$

where $\hat{\sigma}$ is the element of $\text{Sym}(J)$ induced by σ . Moreover, we define $\chi_J(X) = \{\chi_J(x) \mid x \in X_J\}$.

3 Structural results

Some results from [8], in which X -alphabet-almost-simple and X -neighbour-transitive codes with $\delta \geq 3$ are characterised, are stated below. This is our starting point when looking at codes that are X -alphabet-almost-simple and $(X, 2)$ -neighbour-transitive with $\delta \geq 3$, since we then have that C is indeed X -neighbour-transitive. The following definitions are needed first. For a subgroup $T \leq S_q$ define $\text{Diag}_m(T) = \{(h, \dots, h) \in B \mid h \in T\}$.

Definition 3.1. A code C in $H(m, q)$ is *diagonally (X, s) -neighbour-transitive* if C is (X, s) -neighbour-transitive and $X \leq \text{Diag}_m(S_q) \rtimes L$.

Each part of Proposition 3.2 is proved in the relevant citation of [8]. Recall the definitions of: $\pi_J(C)$ and $\chi_J(X)$ (see Section 2.2), the socle $\text{soc}(G)$ and the kernel $K = X \cap B$ for the action of X on M , where $B \cong S_m$ is the base group of $\text{Aut}(\Gamma)$ (see Section 2.1). Note also that G is a *sub-direct* subgroup of a direct product $\prod_{i=1}^n T_i$ of isomorphic groups $T_i \cong T$, where $i \in \{1, \dots, n\}$, if the projection of G in each coordinate is isomorphic to T .

Proposition 3.2. *Suppose C is an X -neighbour-transitive code in $H(m, q)$ with $\delta \geq 3$. Then the following hold:*

- i) *Let \mathcal{J} be an X -invariant partition of M and $J \in \mathcal{J}$ such that $\pi_J(C)$ is not the complete code. Then $\pi_J(C)$ is $\chi_J(X)$ -neighbour-transitive [8, Proposition 3.4]. (Note that the assumption that $\pi_J(C)$ is not the complete code does not appear in [8], but is necessary since the proof assumes that $\pi_J(C)_1$ is non-empty.)*

- ii) Let \mathcal{J} be an X -invariant partition of M and $J \in \mathcal{J}$ such that $\pi_J(C)$ is not the complete code. Then $\pi_J(C)$ has minimum distance at least 2 [8, Corollary 3.7].
- iii) If C is also X -alphabet-almost-simple, then $\text{soc}(K)$ is a sub-direct subgroup of $\prod_{i \in M} \text{soc}(X_i^{Q_i})$ [8, Proposition 5.2].

While the next result is not explicitly stated in [8], it is the basis of the characterisation contained within it.

Proposition 3.3. *Let C be an X -alphabet-almost-simple and X -neighbour-transitive code with $\delta \geq 3$. Then there exists an X -invariant partition \mathcal{J} of M such that for all $J \in \mathcal{J}$ the code $\pi_J(C)$ is equivalent to a diagonally $\chi_J(X)$ -neighbour-transitive with minimum distance $\delta(\pi_J(C)) \geq 2$.*

Proof. Let T be the non-abelian simple socle of the almost-simple 2-transitive group $X_i^{Q_i}$. By Proposition 3.2-(iii), the group $\text{soc}(K)$ is a sub-direct subgroup of $\prod_{i \in M} \text{soc}(X_i^{Q_i})$. Following the discussion after [8, Proposition 5.2], Scott’s Lemma [16, p. 328] can be applied to give a partition \mathcal{J} of M such that $\text{soc}(K) = \prod_{J \in \mathcal{J}} D_J$, where each $D_J \cong \text{Diag}_k(T)$ acts on $\pi_J(V\Gamma)$, for all $J \in \mathcal{J}$, where $k = |J|$. Moreover, by [8, Remark 5.5], \mathcal{J} is X -invariant. By examining $\text{soc}(K)$, it can be shown [8, Section 5] that, up to equivalence, two possibilities occur. Either $\chi_J(X) \leq \text{Diag}_k(S_q) \rtimes S_k$, where $k = |J|$, for all $J \in \mathcal{J}$, or \mathcal{J} can be replaced by a more refined X -invariant partition $\hat{\mathcal{J}}$ of M such that $\chi_{\hat{J}}(X) \leq \text{Diag}_{\hat{k}}(S_q) \rtimes S_{\hat{k}}$, where $\hat{k} = |\hat{J}|$, for all $\hat{J} \in \hat{\mathcal{J}}$.

In either case, it follows from Proposition 3.2-(i) and (ii) that, for all $J \in \mathcal{J}$ or $\hat{\mathcal{J}}$ respectively, $\chi_J(X)$ acts transitively on $\pi_J(C)$ and either $\pi_J(C)$ is the complete code or it is $\chi_J(X)$ -neighbour-transitive with minimum distance at least 2. Since $\chi_J(X)$ is a diagonal subgroup, we deduce that $\pi_J(C)$ is as in the second case, since no diagonal subgroup acts transitively on the complete code. □

Proposition 3.4. *Let C be an $(X, 2)$ -neighbour-transitive code with $\delta \geq 3$ in $H(m, q)$, and suppose \mathcal{J} is an X -invariant partition of M . Then for all $J \in \mathcal{J}$, either;*

- i) $\pi_J(C)$ is the complete code, $\delta(\pi_J(C)) = 1$, and $\chi_J(X)$ is transitive on $\pi_J(C)$;
- ii) $\pi_J(C)$ has covering radius 1, $\delta(\pi_J(C)) = 2$ or 3, and is $(\chi_J(X), 1)$ -neighbour-transitive; or,
- iii) $\pi_J(C)$ is $(\chi_J(X), 2)$ -neighbour-transitive.

Proof. Let $\bar{C} = \pi_J(C)$. The fact that $\chi_J(X)$ is transitive on \bar{C} and \bar{C}_1 , if \bar{C}_1 is non-empty, follows from Proposition 3.2-(i). From this we deduce (i) and (ii). In particular, suppose the covering radius of \bar{C} is at most 1. If the covering radius is 0 then \bar{C} is the complete code, and if the covering radius is 1 then \bar{C} is not the complete code and the minimum distance is at most 3 so, by Proposition 3.2-(ii), the minimum distance is at least 2. Therefore, we need only show that when \bar{C}_2 is non-empty $\chi_J(X)$ is transitive on \bar{C}_2 .

Suppose \bar{C} has covering radius at least 2. Let $\mu, \nu \in \bar{C}_2$. Then there exists $\alpha, \beta \in C$ such that $d(\mu, \pi_J(\alpha)) = d(\nu, \pi_J(\beta)) = 2$. Let $\hat{\nu} \in H(m, q)$ with $\hat{\nu}_u = \nu_u$ for u in J and $\hat{\nu}_v = \alpha_v$ otherwise. Similarly, let $\hat{\mu} \in H(m, q)$ with $\hat{\mu}_u = \mu_u$ for u in J and $\hat{\mu}_v = \beta_v$ otherwise. We claim that $\hat{\nu}, \hat{\mu} \in C_2$. We show this for $\hat{\nu}$ and note that an identical argument holds for $\hat{\mu}$. First, note that $d(\alpha, \hat{\nu}) = 2$ and $\delta \geq 3$, so $\hat{\nu} \notin C$. Suppose $\hat{\nu} \in C_1$. Then

there exists $\alpha' \in C$ such that $d(\hat{\nu}, \alpha') = 1$. We then have $d(\nu, \pi_J(\alpha')) \leq 1$. However, this contradicts $\nu \in \bar{C}_2$. Hence $\hat{\mu}, \hat{\nu} \in C_2$.

As C is $(X, 2)$ -neighbour-transitive, there exists an $x = h\sigma \in X$ mapping $\hat{\nu}$ to $\hat{\mu}$. We claim $x \in X_J$. Suppose $x \notin X_J$. Then, since \mathcal{J} is a system of imprimitivity for the action of X on M , there exists $J' \in \mathcal{J}$ such that $J \neq J'$ and $J'^\sigma = J$. Since $\pi_{J'}(\hat{\nu}) = \pi_{J'}(\alpha)$, this implies that $\pi_J(\hat{\nu}^x) = \pi_J(\alpha^x) \in \bar{C}$ and hence $\pi_J(\hat{\nu}^x) \neq \mu$, which contradicts the fact that $\hat{\nu}^x = \hat{\mu}$. Thus $x \in X_J$ and

$$\nu^{\chi_J(x)} = \pi_J(\hat{\nu})^{\chi_J(x)} = \pi_J(\hat{\nu}^x) = \pi_J(\hat{\mu}) = \mu.$$

□

Proposition 3.5. *Let C be an $(X, 2)$ -neighbour-transitive code in $H(m, q)$ with $\delta \geq 3$, and \mathcal{J} be an X -invariant partition of M . Then, for all $J \in \mathcal{J}$ and $i \in J$,*

1. $\chi_J(X)_i^{Q_i}$ is 2-transitive on Q ; and,
2. for $\alpha \in C$, $\chi_J(X)_{\pi_J(\alpha)}$ is transitive on J .

Proof. As C is X -neighbour-transitive with $\delta \geq 3$, we have that $X_i^{Q_i}$ is 2-transitive, by [7, Proposition 2.7], and X^M is transitive, by [7, Proposition 2.5]. One then deduces that $X_i^{Q_i}$ is 2-transitive for all i . Now, because \mathcal{J} is an X -invariant partition, it follows that $X_i = (X_J)_i$ for all $i \in J$. This in turn implies that $\chi_J(X)_i = \chi_J(X_i)$. It is now straight forward to show that $\chi_J(X_i)^{Q_i} = X_i^{Q_i}$.

Now, since X_α is transitive on M and \mathcal{J} is an X -invariant partition of M , it follows that $(X_\alpha)_J$ is transitive on J . Thus $\chi_J(X_\alpha) \leq \chi_J(X)_{\pi(\alpha)}$ is transitive on J . □

The previous two propositions suggest a study of codes that are $(X, 2)$ -neighbour-transitive, have minimum distance $\delta \geq 2$, and where X acts primitively on M . An answer to the following questions would provide us with the building blocks for $(X, 2)$ -neighbour-transitive codes with $\delta \geq 3$.

Question 3.6. Can we classify all $(X, 2)$ -neighbour-transitive codes with $\delta \geq 2$ such that X^M is primitive and $X_i^{Q_i}$ is 2-transitive?

Question 3.7. Can we classify all $(X, 1)$ -neighbour-transitive codes with $\delta = 2$ or 3 and $\rho = 1$ such that X^M is primitive and $X_i^{Q_i}$ is 2-transitive?

Let C be a code and $X \leq \text{Aut}(C)$. If X acts faithfully on M , that is $K = X \cap B = 1$, we say C is X -entry-faithful. If $K \neq 1$, X^M is transitive on M and $X_i^{Q_i}$ is affine ($X_i^{Q_i} \leq \text{AGL}_d(p)$ for some integer d and prime p) we say C is X -alphabet-affine. Questions 3.6 and 3.7 can be further broken down into X -entry-faithful and non-trivial kernel cases, that is, X -alphabet-affine and X -alphabet-almost-simple (see Section 2.1 for the definition of X -alphabet-almost-simple). By the main result of this paper, the outstanding cases of Question 3.6 are X -alphabet-almost-simple and $(X, 2)$ -neighbour-transitive with $\delta = 2$, and X -alphabet-affine and $(X, 2)$ -neighbour-transitive, where X^M is primitive and $X_i^{Q_i}$ is 2-transitive.

Given Proposition 3.3, a third question is the following.

Question 3.8. Can we construct $(X, 2)$ -neighbour-transitive codes with $\delta \geq 3$ by taking copies of $(X, 1)$ -neighbour-transitive codes with $\delta = 2$ or 3 and $\rho = 1$.

4 Examples

We begin this section by considering some examples of codes which have properties relating to the results of the previous section. We first introduce the operators Prod and Rep which allow the construction of new codes from old ones. For an arbitrary code C in $H(m, q)$ we define $\text{Prod}(C, \ell)$ and $\text{Rep}_\ell(C)$ in $H(m\ell, q)$ as

$$\text{Prod}(C, \ell) = \{(\alpha_1, \dots, \alpha_\ell) \mid \alpha_i \in C\},$$

and

$$\text{Rep}_\ell(C) = \{(\alpha, \dots, \alpha) \mid \alpha \in C\}.$$

The *repetition code* $\text{Rep}(m, q)$ in $H(m, q)$ is the set of all vertices (a, \dots, a) consisting of a single element $a \in Q$ repeated m times.

The next two examples present codes which are both X -alphabet-almost-simple and X -completely transitive, though the second example has minimum distance $\delta = 2$.

Example 4.1. Let $C = \text{Rep}(3, q)$, where $q \geq 5$, and $X = \text{Diag}_3(S_q) \rtimes S_3$, as in [11, Example 3.1]. Now,

$$C_1 = \{(a, a, b), (a, b, a), (b, a, a) \mid a, b \in Q; a \neq b\},$$

and

$$C_2 = \{(a, b, c) \mid a, b, c \in Q; a \neq b \neq c \neq a\}.$$

Since S_q acts 3-transitively on Q and S_3 acts transitively on M , it follows that X acts transitively on C, C_1 and C_2 . Thus C is $(X, 2)$ -neighbour-transitive and X -completely transitive, since C has covering radius $\rho = 2$. Also, $X_i^{Q_i} \cong S_q$ is almost-simple, since $q \geq 5$, and $X^M \cong S_3$ is transitive on M . Hence C is X -alphabet-almost-simple and X -completely transitive.

Example 4.2. Let $q \geq 5, \ell \geq 2, C = \text{Prod}(\text{Rep}(2, q), \ell)$ and $X = (\text{Diag}_2(S_q))^\ell \rtimes U$, where $\text{Diag}_2(S_q)$ is a subgroup of the base group of $\text{Aut}(H(2, q))$ and $U = S_2 \wr S_\ell = S_2^\ell \rtimes S_\ell$ is a subgroup of the top group of $\text{Aut}(H(2\ell, q))$. Let $\mathcal{J} = \{J_1, \dots, J_\ell\}$, with $J_i = \{2i - 1, 2i\}$, be the partition of M preserved by U . Note that $\delta = 2$. Let $R \subseteq \{1, \dots, \ell\}$ of size s , and $\nu \in H(m, q)$ be such that $\pi_{J_i}(\nu) = (a, b)$, where $a \neq b$ for all $i \in R$, and $a = b$ for all $i \notin R$. Any codeword β is at least distance s from ν , since $d(\pi_{J_i}(\nu), \pi_{J_i}(\beta)) \geq 1$ for each $i \in R$. Also, there exists some codeword α with $\pi_{J_i}(\alpha) = (a, a)$ whenever $\pi_{J_i}(\nu) = (a, b)$ for $i \in \{1, \dots, \ell\}$, and hence $d(\alpha, \nu) = s$. So $\nu \in C_s$. Any vertex ν of $H(2\ell, q)$ can be expressed in this way, for some R , since $\pi_{J_i}(\nu) = (a, b)$ has either $a = b$ or $a \neq b$. Thus, for each s, C_s consists of all such vertices ν where $|R| = s$. It also follows from this that $\rho = \ell$.

Let $\nu \in C_s$, with R as above. Let $x = (h_{J_1}, \dots, h_{J_\ell})\sigma \in X$ where $h_{J_i} \in \text{Diag}_2(S_q)$ such that $\pi_{J_i}(\nu)^{h_{J_i}} = (1, 2)$, for $i \in R$, and $\pi_{J_i}(\nu)^{h_{J_i}} = (1, 1)$, for all $i \notin R$. Moreover, since S_ℓ is ℓ -transitive, there exists $\sigma \in S_\ell \leq S_2 \wr S_\ell$ mapping $\{J_{i_1}, \dots, J_{i_s}\}$ to $\{J_1, \dots, J_s\}$ (where $R = \{i_1, \dots, i_s\}$), whilst preserving order within each J_i . Then $\nu^x = \gamma \in C_s$, where $\pi_{J_i}(\gamma) = (1, 2)$ for all $i \in \{1, \dots, s\}$ and $\pi_{J_i}(\gamma) = (1, 1)$ for all $i \notin \{s + 1, \dots, \ell\}$. Since we can map any such ν to γ, X is transitive on C_s for each $s \in \{1, \dots, \ell\}$. Hence C is X -completely transitive, and in particular $(X, 2)$ -neighbour-transitive for $\ell \geq 2$. Since $X_i^{Q_i} \cong S_q$ and $X^M \cong S_2 \wr S_\ell$ is transitive on M, C is X -alphabet-almost-simple.

Lemma 4.3. *Suppose C is an $(X, 2)$ -neighbour-transitive code in $H(m, q)$, with $q \geq 3$, and \mathcal{J} is an X -invariant partition of M , such that $\pi_J(C) = \text{Rep}(k, q)$, for all $J \in \mathcal{J}$ where $k = |J|$. Then either $\delta = k = 2$, or \mathcal{J} is a trivial partition.*

Proof. Let $x = (h_1, \dots, h_m)\sigma \in X$ and $J \in \mathcal{J}$. By the hypothesis it follows that for all $a \in Q$, there exists $\alpha \in C$ such that $\pi_J(\alpha) = (a, \dots, a)$. Suppose $J^\sigma = J' \in \mathcal{J}$. Then $\pi_{J'}(\alpha^x) = (a^{h_{i_1}}, \dots, a^{h_{i_k}})\sigma = (b, \dots, b)$ for some $b \in Q$, that is, $a^{h_{i_s}} = a^{h_{i_t}}$ for all $i_s, i_t \in J$. In particular $\chi_J(x\sigma^{-1}) = (h, \dots, h)$ for some $h \in S_q$, and $X \leq \text{Diag}_k(S_q)\wr U$, where U is the stabiliser of \mathcal{J} in the top group.

Suppose that the partition \mathcal{J} is non-trivial, so that $k, \ell \geq 2$. Since C is a subset of $\text{Prod}(\text{Rep}(k, q), \ell)$, which has minimum distance k , it follows that $\delta \geq k \geq 2$.

Suppose $\delta \geq 3$. As C is a subset of $\text{Prod}(\text{Rep}(k, q), \ell)$ we can replace C by an equivalent code contained in $\text{Prod}(\text{Rep}(k, q), \ell)$ containing $\alpha = (1, \dots, 1)$ and such that

$$\mathcal{J} = \{\{1, \dots, k\}, \{k + 1, \dots, 2k\}, \dots, \{m - k + 1, \dots, m\}\}.$$

Consider,

$$\begin{aligned} \mu &= (2, 3, 1, 1, \dots, 1, 1, 1, 1, \dots, 1, \dots, 1, \dots, 1) \quad \text{and} \\ \nu &= (\underbrace{2, 1, 1, 1, \dots, 1}_{k \text{ entries}}, \underbrace{2, 1, 1, \dots, 1}_{k \text{ entries}}, \dots, \underbrace{1, \dots, 1}_{k \text{ entries}}). \end{aligned}$$

If $k = 2$, then we claim $\mu \in C_2$. Any vertex $\beta \in \text{Prod}(\text{Rep}(2, q), \ell) \supseteq C$ with $d(\mu, \beta) = 1$ is of the form $\gamma = (a, a, 1, \dots, 1)$, where $a = 2$ or 3 . However, no such γ is an element of C , since each is distance 2 from α . If $k \geq 3$ then $\mu \in C_2$ since $d(\alpha, \mu) = 2$ and there is no closer codeword as $\pi_{J_1}(\mu) \in \pi_{J_1}(C)_2$. In both cases $\nu \in C_2$ since $d(\alpha, \nu) = 2$ and no codeword is closer, as $\pi_{J_i}(\nu) \in \pi_{J_i}(C)_1$ for $i = 1, 2$. Let $x = (h_1, \dots, h_m)\sigma \in X$ such that $\mu^x = \nu$. We reach a contradiction here, since $h_1 = h_2 = \dots = h_k = h$ cannot, assuming $k \geq 3$, map the set $\{1, 2, 3\}$ to either of the sets $\{1, 2\}$ or $\{1\}$. In the case $k = 2$, in at least one block we must map the set $\{1\}$ to $\{1, 2\}$, which is not possible. Hence $2 \geq \delta \geq k \geq 2$. □

Suppose \mathcal{J} is a system of imprimitivity for the action of X on M and C is an X -neighbour-transitive code, with $\delta \geq 3$. The next example shows that it is possible that the projection of C onto each block of \mathcal{J} gives the complete code, though this is not the system of imprimitivity of interest to us in Proposition 3.3.

Example 4.4. Let $\bar{C} = \text{Prod}(C, \ell)$ be a code in $\Gamma = H(m, q)$, where $m = k\ell$ and C is an X -neighbour-transitive code in $H(k, q)$ where $X \cap B$ is transitive on C and $\delta \geq 3$. Let $\bar{X} = \langle (X \cap B)^\ell, \text{Diag}_\ell(X), S_\ell \rangle$ preserve the partition

$$\mathcal{J} = \{\{1, \dots, k\}, \dots, \{m - k + 1, \dots, m\}\} = \{J_1, \dots, J_\ell\},$$

of M , where $\chi_J((X \cap B)^\ell) = X \cap B$ and $\chi_J(\text{Diag}_\ell(X)) = X$ for all $J \in \mathcal{J}$, and S_ℓ acts as pure permutations by permuting the blocks of \mathcal{J} whilst preserving the order of entries within a given block. It follows that we preserve two \bar{X} -invariant partitions. These being \mathcal{J} and \mathcal{J}' , where \mathcal{J}' is attained by taking the corresponding entries, by order, from each copy of C to form each block:

$$\mathcal{J}' = \{\{1, k + 1, \dots, m - k + 1\}, \dots, \{\ell, k + \ell, \dots, m\}\}.$$

Given any $\alpha = (\alpha_1, \dots, \alpha_\ell) \in \bar{C}$, $\alpha_i \in C$, and $\beta = (\beta_1, \dots, \beta_\ell) \in \bar{C}$, $\beta_i \in C$ there exists an $x \in (X \cap B)^\ell$ mapping α to β since $X \cap B$ is transitive on C . Hence \bar{X} is transitive on \bar{C} . Given any two neighbours $\mu, \nu \in \Gamma_1(\alpha)$, where μ, ν differ from α in the respective blocks J_i and J_j , we can map J_j to J_i via some element $\sigma \in S_\ell$. Then, since X_{α_i} is transitive on $\Gamma_1(\alpha_i)$, there exists an element $x \in \text{Diag}_\ell(X)$ such that $\pi_{J_i}(\nu^{\sigma x}) = \pi_{J_i}(\mu)$. We can then map $\nu^{\sigma x}$ to μ via some element $h \in (X \cap B)^\ell$, where $\chi_{J_i}(h) = 1$, since each $\pi_{J_t}(\nu^{\sigma x})$ and $\pi_{J_t}(\mu)$ are elements of C for $t \neq i$ and $X \cap B$ is transitive on C . Hence $\sigma x h$ maps ν to μ and \bar{X} is transitive on \bar{C}_1 .

When we consider the projection $\pi_J(\bar{C})$ for any $J \in \mathcal{J}'$ we are left with the complete code. To see this, consider that for $(\alpha_1, \dots, \alpha_\ell) \in \bar{C}$, $\alpha_i \in C$, we may choose an arbitrary element of C as α_i for each i . Since $X_i^{Q_i}$ is 2-transitive on Q_i , each element appears in the first entry for some codeword. Thus, as $\pi_J((\alpha_1, \dots, \alpha_\ell))$ when $J = \{1, k + 1, \dots, m - k + 1\}$ is the first entry of each α_i , we have that $\pi_J(\bar{C})$ is the complete code.

5 Alphabet-almost-simple $(X, 2)$ -neighbour-transitive codes

Before we prove the final results we define the codes used in this section, which first requires the following definition.

Definition 5.1. Define the *composition* of a vertex $\alpha \in H(m, q)$ to be the set

$$Q(\alpha) = \{(a_1, p_1), \dots, (a_q, p_q)\},$$

where p_i is the number of entries of α which take the value $a_i \in Q$. For $\alpha \in H(m, q)$ define the set

$$\text{Num}(\alpha) = \{(p_1, s_1), \dots, (p_j, s_j)\},$$

where (p_i, s_i) means that s_i distinct elements of Q appear precisely p_i times in α .

Definition 5.2. We define the following codes:

1. $\text{Inj}(m, q)$, where $m < q$, is the set of all vertices $\alpha \in H(m, q)$ such that $\text{Num}(\alpha) = \{(1, m)\}$;
2. for m odd, $W(\lfloor m/2 \rfloor, 2)$ is the set of vertices in $\alpha \in H(m, 2)$ such that $\text{Num}(\alpha) = \{(m + 1)/2, 1), (m - 1)/2, 1)\}$; and,
3. $\text{All}(pq, q)$, with $p q = m$, is the set of all vertices $\alpha \in H(m, q)$ such that $\text{Num}(\alpha) = \{(p, q)\}$.

More information on these codes is available in [9, Definition 2]. The following lemma is [9, Lemma 4].

Lemma 5.3. Let α be a vertex in $H(m, q)$. Then $\text{Num}(\alpha)$ is preserved by $\text{Diag}_m(S_q) \rtimes L$.

The last result, in combination with the classification of diagonally neighbour-transitive codes [9, Theorem 4.3], allows us to prove the next result.

Proposition 5.4. Let C be a diagonally $(X, 2)$ -neighbour-transitive code in $H(m, q)$. Then one of the following holds:

1. $q = 2$ and $C = \{(a, \dots, a)\}$;
2. $m = 3$ or $q = 2$, and $C = \text{Rep}(m, q)$;

3. $C = \text{Inj}(3, q)$;
4. m is odd and $C = W([m/2], 2)$; or,
5. $q = 2$ or $q = m = 3$, and there exists some p such that $m = pq$ and C is a subset of $\text{All}(pq, q)$.

Proof. From [9, Theorem 4.3], we have that a diagonally neighbour-transitive code C is one of: $\{(a, \dots, a)\}$ for some $a \in Q$, $\text{Rep}(m, q)$, $\text{Inj}(m, q)$ with $m < q$, $W([m/2], 2)$ with m odd, or there exists a p such that $m = pq$ and C is a subset of $\text{All}(pq, q)$. Here we consider $m \geq 2$, since if $m = 1$ then C_2 is empty, so C is not $(X, 2)$ -neighbour-transitive. Also to prove some C is $(X, 2)$ -neighbour-transitive, we need only find some $X \leq \text{Aut}(C)$ such that $X \leq \text{Diag}_m(S_q) \rtimes L$ and X is transitive on C_2 , since C is already X -neighbour-transitive, for some X , by [9, Theorem 4.3].

First, if $C = \text{Inj}(2, q)$ then C_2 is empty. Thus, C is not $(X, 2)$ -neighbour-transitive. Table 1 lists the remaining cases which are not 2-neighbour-transitive. The second and third columns give a pair $\mu, \nu \in C_2$ such that $\text{Num}(\mu) \neq \text{Num}(\nu)$. Hence, by Lemma 5.3, X is not transitive on C_2 . It can be deduced from $\text{Num}(\mu), \text{Num}(\nu)$ that $\mu, \nu \in C_2$, since this makes it clear that we must change μ, ν in at least two entries to get a vertex in C . Note that we let $\alpha = (1, 2, 3, \dots, q) \in H(q, q)$ and in the second last and last rows we assume $\alpha \in C$ and $(\alpha, \dots, \alpha) \in C$, respectively, and observe for the last row $\hat{\mu} = (1, 1, 1, 4, 5, \dots, q)$, $\hat{\nu} = (1, 1, 3, 4, 5, \dots, q)$ are in $\Gamma_2(\alpha)$.

C Conditions	$\mu \in C_2$ Num(μ)	$\nu \in C_2$ Num(ν)
$\{(a, \dots, a)\}$ $q \geq 3$	(b, b, a, \dots, a) $\{(m - 2, 1), (2, 1)\}$	(b, c, a, \dots, a) $\{(m - 2, 1), (1, 2)\}$
$\text{Rep}(m, q)$ $m > q \geq 3$	$(2, 2, 1, \dots, 1)$ $\{(m - 2, 1), (2, 1)\}$	$(2, 3, 1, \dots, 1)$ $\{(m - 2, 1), (1, 2)\}$
$\text{Inj}(m, q)$ $m \geq 4$	$(1, 1, 1, 4, 5, \dots, m)$ $\{(3, 1), (1, m - 3)\}$	$(1, 1, 3, 3, 5, 6, \dots, m)$ $\{(2, 2), (1, m - 4)\}$
$\subseteq \text{All}(q, q)$ $q \geq 4$	$(1, 1, 1, 4, 5, \dots, q)$ $\{(3, 1), (1, q - 3)\}$	$(1, 1, 3, 3, 5, 6, \dots, q)$ $\{(2, 2), (1, q - 4)\}$
$\subseteq \text{All}(pq, q)$ $q > p \geq 2$	$(\hat{\mu}, \alpha, \dots, \alpha)$ $\{(p - 1, 2), (p, q - 3), (p + 2, 1)\}$	$(\hat{\nu}, \hat{\nu}, \alpha, \dots, \alpha)$ $\{(p - 2, 1), (p, q - 2), (p + 2, 1)\}$

Table 1: Diagonally neighbour-transitive codes C which are not diagonally 2-neighbour-transitive, and elements of C_2 which illustrate this. Note: $\hat{\mu} = (1, 1, 1, 4, 5, \dots, q)$, $\hat{\nu} = (1, 1, 3, 4, 5, \dots, q)$ and $\alpha = (1, 2, 3, \dots, q)$.

Now we prove the result for the cases which are 2-neighbour-transitive. Suppose $C = \{(a, \dots, a)\}$ for some $a \in Q$. Let $q = 2$ and $Q = \{0, 1\}$. Then $L = S_m = \text{Aut}(C)$. Without loss of generality, let $a = 0$ so that C_2 is the set of weight two vertices. Since L is transitive on the sets of weight 2 and weight 1 vertices, it follows C is diagonally $(X, 2)$ -neighbour-transitive. Let $C = \text{Rep}(m, q)$. It follows from Example 4.1 that $\text{Rep}(3, q)$ is $(\text{Diag}_3(S_q) \rtimes S_3, 2)$ -neighbour-transitive. If $q = 2$ then $\text{Aut}(C) \cong \text{Diag}_m(S_2) \rtimes S_m$ and C is completely transitive [11, Example 3.1]. Consider $C = \text{Inj}(m, q)$ with $3 = m < q$ and $q \geq 4$. If $\nu \in C_2$ then $\nu_1 = \nu_2 = \nu_3$, since otherwise $\nu \in C$ or C_1 . Since $\text{Diag}_m(S_q) \leq \text{Aut}(C)$, we are transitive on C_2 . Suppose $C = W([m/2], 2)$ and m is odd. Then by [9, Corollary 3.4] C is $\text{Diag}(S_2) \rtimes S_m$ -completely transitive. Finally, suppose C is a subset

of $\text{All}(pq, q)$ for some p such that $m = pq$. Let $p \geq 2$, $q = 2$ and $C = \text{All}(2p, 2)$. Then C_2 is the set of all weight $p \pm 2$ vertices, which $\text{Diag}_2(S_2) \rtimes S_m \leq \text{Aut}(C)$ is transitive on. Let $p = 1$, $q = 3$ and $C = \text{All}(3, 3)$. Then $C_2 = \text{Rep}(3, q)$ and is $\text{Aut}(C)$ -completely transitive by Example 4.1. \square

With our classification of diagonally $(X, 2)$ -neighbour-transitive codes from the previous result, Propositions 3.3 and 3.4 mean we are now in a position to prove the main theorem.

Proof of Theorem 1.1. Suppose C is an X -alphabet-almost-simple and $(X, 2)$ -neighbour-transitive code with $\delta \geq 3$ such that $X \cap B \neq 1$. By Proposition 3.3, there exists an X -invariant partition $\mathcal{J} = \{J_1, \dots, J_\ell\}$, for some ℓ , for the action of X on M . Moreover, $\pi_{J_i}(C)$ has minimum distance at least 2 and is diagonally $\chi_{J_i}(X)$ -neighbour-transitive. By Proposition 3.4, either $\pi_{J_i}(C)$ has covering radius $\rho \leq 1$, or $\pi_{J_i}(C)$ is also $(\chi_{J_i}(X), 2)$ -neighbour-transitive. Note $\rho \neq 0$, that is, $\pi_{J_i}(C)$ is not the complete code, since $\pi_{J_i}(C)$ has minimum distance at least 2.

Suppose $\pi_{J_i}(C)$ has covering radius $\rho \geq 2$. Since $X_i^{Q_i}$ is almost-simple, it follows that $q \geq 5$. By Proposition 5.4, the only diagonally 2-neighbour-transitive code with $q \geq 5$ and $\delta \geq 2$ is $\text{Rep}(3, q)$ for $q \geq 5$ (note that $\delta = 1$ for $\text{Inj}(3, q)$). Then Lemma 4.3 implies \mathcal{J} is a trivial partition. Since $|J_i| = k = 3 > 1$, it follows that $\ell = 1$, $k = m$, and $C = \text{Rep}(3, q)$.

Suppose $\pi_{J_i}(C)$ has covering radius $\rho = 1$. Now, by [9, Thm. 4 and Cor. 2], the only diagonally neighbour-transitive code with $\delta \geq 2$ and $\rho = 1$ is $\text{Rep}(2, q)$. If $l = 1$ then we have $\delta = 2$, a contradiction. Suppose $l \geq 2$. Then Lemma 4.3 implies $\delta = 2$, a contradiction. \square

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