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Large Graph Clustering Using DCT-Based Graph Clustering

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Abstract—With the proliferation of the World Wide Web, graph structures have arisen on social network/media sites. Such graphs usually number several million nodes, i.e., they can be characterized as Big Data. Graph clustering is an important analysis tool for other graph related tasks, such as compression, community discovery and recommendation systems, to name a few. We propose a novel extension to a graph clustering algorithm, that attempts to cluster a graph, through the optimization of selected terms of the graph weight/adjacency matrix Discrete Cosine Transform.

I. INTRODUCTION

Graph clustering, i.e., grouping graph nodes that are closely interconnected has an important role in graph analysis applications. In this paper, we will present a genetic algorithm for graph clustering that is designed to be fast and whose memory requirements are linear with respect to the number of both graph nodes and graph edges.

There has been an abundance of works that deal with the clustering of graph nodes. For an in-depth survey on graph clustering, the interested reader can refer to [1]. One of the more theoretically elegant types of approaches is using the graph spectrum [2] to perform the clustering. Such methods are referred to as spectral graph clustering methods, a popular example of which is the Normalized Cut, or NCUT, algorithm [3].

We base the method proposed in this paper on a novel approach to graph node clustering presented in [4], which uses optimization on selected terms of the Discrete Cosine Transform (DCT) [5] of the graph adjacency matrix to perform the clustering. We will attempt to solve the parameterization issues that were observed in that work, by a) constructing a padded graph based on the input graph, b) defining a different DCT-related objective function to optimize, c) restricting the way optimization can be performed and d) using the new method to recursively split the clusters into subclusters.

The paper is organized as follows: Section II briefly describes the initial idea presented in [4], section III details the modifications that are proposed on a theoretical level, while section IV describes how the computations can be significantly sped up. Section V provides the results of the comparative evaluation experiments and Section VI concludes the paper.

II. DCT-BASED GRAPH CLUSTERING

As proposed in [4], the main idea behind using the DCT of the graph adjacency matrix to perform clustering on a graph, stems from a simple observation. When the nodes of each cluster appear in consecutive places in the order used to construct the adjacency matrix, then large blocks of edges appear along the image diagonal. This is illustrated in Figure 1.

In the DCT domain, large, square blocks along the diagonal correspond to large values in the diagonal terms of the DCT, as can be seen in Figure 2. Swapping the order of nodes in the graph adjacency matrix causes the DCT term values to fluctuate. It is, therefore, possible to bring nodes of the same cluster together in the adjacency matrix by maximizing the sum of the first $k$ terms of the DCT diagonal. This can be achieved by using Simulated Annealing (SA) [4].

The issue that arises, however, is the determination of the parameter $k$. In general, $k$ should be close to the number of clusters. As Figure 3 shows, using the wrong values for this parameter can cause problems. In this example, there are 3 clusters in the input graph. If the number of clusters is underestimated, then the two smaller clusters are merged together (Figure 3b). If the number of clusters is overestimated, then the biggest cluster is split (Figure 3d). This can be useful, if we want the graph to be clustered into a specific number of clusters. However, in the general case, we must correctly guess the number of clusters, otherwise the method will underperform. Furthermore, depending on the size difference between the smallest and biggest cluster, it may be impossible to choose a value for $k$ that will not split the biggest cluster or merge the smaller one into another cluster. It is this issue that we address in this paper.
This section details our proposed modifications to [4] that aim to de-parameterize the method and improve clustering performance. The optimization method remains Simulated Annealing.

We will first describe the changes to the input graph adjacency matrix of $N$ nodes. We pad the matrix with dummy nodes to obtain a padded graph with adjacency matrix $B$ as follows:

- **Group 1**: nodes in matrix places 1-$N$ are dummy nodes, not connected to any other node.
- **Group 2**: nodes in matrix places $(N+1)$-$(3N/2)$ are auxiliary nodes, connected to every node in groups 2, 3 and 4 with weight 0.5.
- **Group 3**: nodes in matrix places $(3N/2+1)$-$(5N/2)$ are the real nodes of the input graph, connected in the same way with each other and also groups 2 and 4 (weight 0.5).
- **Group 4**: nodes in matrix places $(5N/2+1)$-3$N$ are auxiliary nodes, connected to every node in groups 2, 3 and 4 with weight 0.5.
- **Group 5**: nodes in matrix places $(3N+1)$-8$N$ are dummy nodes, not connected to any other node.

This creates an $8N \times 8N$ matrix, as illustrated in Figure 4.

Now we will describe the restrictions we place on node swapping. Instead of allowing any two nodes to swap places in the adjacency weight matrix, we attempt to optimize the objective function $f(B)$ by swapping nodes in the following way:

- Select a random node with position $i$ in the range of $N$-3$N$.
- Swap the selected node with the node in position $5N+i$.
- Calculate the change in the objective function $f(B)$.
- Accept the swap (transition) with the appropriate probability, as per SA operation.

In order to visualize what happens with these swap, please refer to Figure 5. When the optimization starts, all the non-zero weights are concentrated in the square marked $C_1$. When the first node is swapped from $C_1$, some of its weights move to squares $C_2$, $I_1$ and $I_2$. As the optimization continues weights move between these squares. Note that squares $I_1$ and $I_2$ contain the weights of the edges connecting the nodes in $C_1$ with the nodes in $C_2$. Suppose that the graph has 2 clusters. It is easy to see that, if every node of one cluster was in $C_1$ and every node of the other cluster was in $C_2$, then the weights in $C_1$ and $C_2$ would be very strong, as they correspond to internal edges between nodes of the same cluster. Respectively, weights in $I_1$ and $I_2$ would be fairly weaker, as they correspond to inter-cluster edges between the 2 clusters. If there are more than 2 clusters in the graph, it suffices that $I_1$ and $I_2$ only contain inter-cluster edges, since we can recursively cluster the nodes of $C_1$ and $C_2$ in the same way, until a stopping criterion is met.

The task, at this point, is to devise an objective function that a) minimizes the inter-cluster weights in $I_1$ and $I_2$ and
employ an exponential function $e^{-\alpha q_2}$ so that, when $q_2$ is high, $e^{-\alpha q_2} \gg q_1$. This forces the optimization to move at least some nodes in $C_2$. Additionally, once $q_2$ reaches a fraction of its initial value (dictated by $\alpha$), then $e^{-\alpha q_2}$ becomes almost a constant and will not force exactly even graph splits. The objective function for the optimization is, therefore, defined to be:

$$f(B) = q_1 + Ne^{-\alpha q_2}. $$

When the minimization is over, the input graph has been partitioned into 2 subgraphs. We can apply the same method again to each subgraph and keep splitting the resulting graphs, until a subgraph is considered too dense to be further split.

IV. Fast implementation

In our explanation of the algorithm, we used the DCT of a padded graph matrix. Using such a matrix in practice requires $O(n^2)$ memory and such an implementation would be infeasible. Additionally, recomputing the DCT is extremely inefficient. It is, in fact, possible to perform all the required computations through a node adjacency list. This is even more useful, when the graph is rather sparse.

Let us consider $DCT_B(4, 2)$, as all the other terms are similar with this one. Note that

$$DCT_B(4, 2) = \sum_{i=1}^{n} \sum_{j=1}^{n} b_{ij} \cos(\frac{\pi}{n} (j + \frac{1}{2}) 2) \cos(\frac{\pi}{n} (i + \frac{1}{2}) 4).$$

The DCT is computed once at the start of the process. We maintain the current position of each node in a global vector and its adjacency list in its own adjacency vectors. The contribution of graph node $k$ that is currently in position $i$ to $DCT_B(4, 2)$ is

$$\sum_{j=1}^{n} b_{ij} \cos(\frac{\pi}{n} (j + \frac{1}{2}) 2) \cos(\frac{\pi}{n} (i + \frac{1}{2}) 4).$$

Moving a node from one group to the other essentially changes the position of the node from $i$ to $i'$. We can calculate the current contribution, by going through its adjacency list vector and retrieving the current positions of its neighbors from the global vector, thus determining the value for $b_{ij}$. Nodes that are not connected to node $k$ yield $b_{ij} = 0$ and are not considered in the calculation.

After computing the contribution, we can subtract it from $DCT_B(4, 2)$. We can then move the node by updating its position in the global position vector, calculate the node’s contribution in the new position and add it back into $DCT_B(4, 2)$. In this way, the changes in the objective function can be calculated in a minimal way.

In this implementation, if $N$ is the number of nodes and $M$ is the number of edges, then the memory requirements for each node $i$ are storing its adjacency lists, which has a size of $d_i$, where $d_i$ is the node’s degree, plus its place in the padded matrix. Note that $\sum_i d_i = M$ and, therefore, the total memory requirements are $\mathcal{O}(M + N)$.

V. Experiments

In order to compare our novel clustering method with NCUT, we ran our experiments on the same data as the facial image clustering method, presented in [6], which uses NCUT. The facial images were extracted from the movies included in the database described in [7]. The weight matrix $A$ was constructed using a Normalized Mutual Information measure, as described in [8]. The matrix was then normalized as follows: $\frac{1}{\max(A)}(A - \min(A))$. The clustering was achieved through the optimization process detailed in Section III. The criterion for stopping the recursive cluster splitting process was checking whether $\max(A) - \min(A) \leq 0.7$. Results for the proposed method and the method in [6] are shown in Table I. Overviewing the results, we conclude that the performance of the proposed method is comparable to that of [6].

In order to evaluate our approach on a large graph, we selected the youtube dataset from the Stanford Network Analysis Project [9]. The dataset is provided with community ground truth. We selected to only include communities, which have more than 100 members. This resulted in 105 communities.
The extracted graph of users that belong to these communities contains 17126 nodes and 212136 edges. The ground truth labels for each node were determined, according the largest community the node belonged to. The comparison between the labels provided by the proposed approach and the ground truth was carried out using the F-measure [10]. The runtime was 33 seconds and the recorded F-measure was 0.1482, which is consistent with the NMI performance, another measure closely related to F-measure, of other state of the art clustering algorithms on Big Data graph clustering [11].

VI. CONCLUSIONS

In this paper we have presented a novel improvement on an original approach to graph clustering and used it to perform facial image clustering and youtube community clustering. By properly constructing a padded graph and only allowing very specific node swaps we were able to define an objective function, whose minimization leads to splitting the graphs in such a way that clusters remain mostly together. Experimental results indicate that this approach performs very comparably with the spectral graph clustering approach in [6]. Also note that the approach in [6] requires providing different values for 2 parameters in order to perform this well, while the method presented here does not require changing values for any parameters. Additionally, we used our approach to perform fast and efficient clustering on a large graph.

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REFERENCES


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